

Automatic Code Generation of Distributed Parallel Tasks

CSE 2016

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Motivation

Scientific Program

Signal Processing

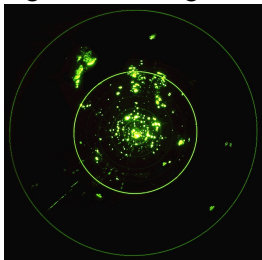
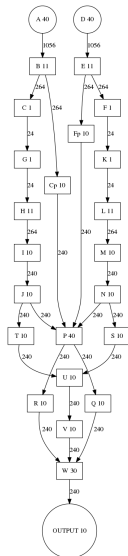


Image Processing



Tools

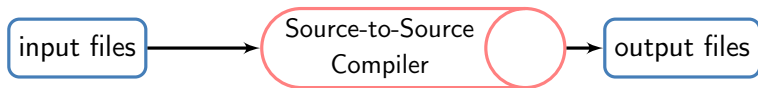
- Automatic task parallelization (OpenModelica¹)
- Automatic distributed parallelization (Pluto+²)
- *Black box*

But no automatic distributed parallelization task tool

¹Mahder Gebremedhin and Peter Fritzon. “Automatic Task Based Analysis and Parallelization in the Context of Equation Based Languages”. In: *Proceedings of the 6th International Workshop on Equation-Based Object-Oriented Modeling Languages and Tools*. EOOLT '14. Berlin, Germany: ACM, 2014, pp. 49–52.

²Uday Bondhugula. “Compiling Affine Loop Nests for Distributed-memory Parallel Architectures”. In: *Proceedings of the International Conference on High Performance Computing, Networking, Storage and Analysis*. SC '13. Denver, Colorado: ACM, 2013.

Source-to-Source Transformations



- Fortran code
- C code

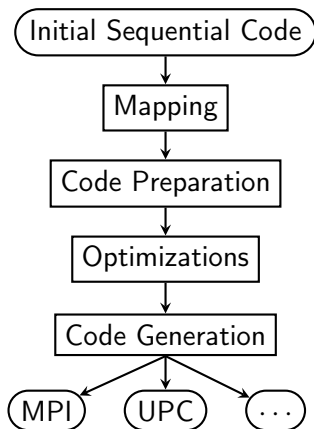
-
- *Scientific Program*
 - *Image Processing*

- Static analyses
- Instrumentation/
Dynamic analyses
- Transformations
- Source code generation
- Code modelling
- Prettyprint

- Fortran code
- C code

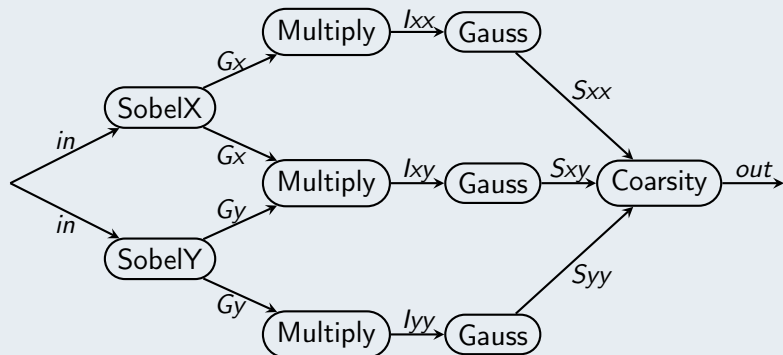
-
- *Distributed Parallel Code*

Compilation Process



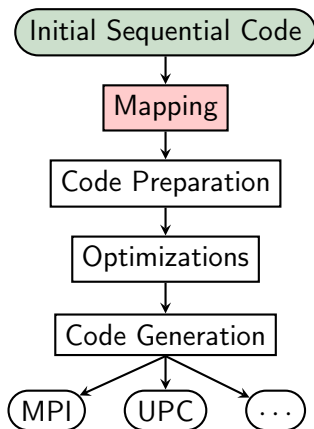
Case Example

Harris&Stephens algorithm³



³Chris Harris and Mike Stephens. "A combined corner and edge detector". In: *In Proc. of Fourth Alvey Vision Conference*. 1988, pp. 147–151

Compilation Process



Can be done

- Automatically with a task scheduler⁴
- Manually

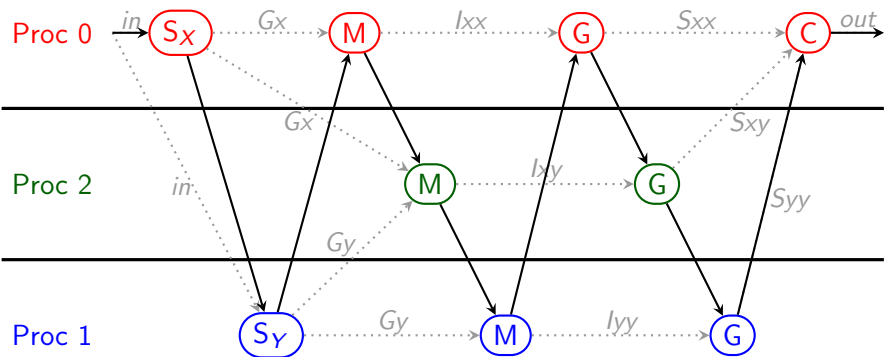
Pragma directive

- New pragma distributed
- On sequence of instructions, loop, test, etc.
- Not inside loop or condition
- `on_cluster` to define the process to use
- No data dependence information needed ⁵

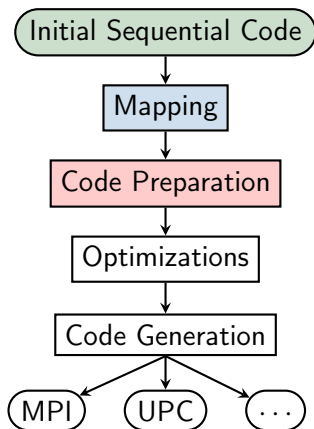
⁴Dounia Khaldi, Pierre Jouvelot, and Corinne Ancourt. “Parallelizing with BDSC, a Resource-constrained Scheduling Algorithm for Shared and Distributed Memory Systems”. In: *Parallel Comput.* 41.C (Jan. 2015), pp. 66–89

⁵Martin Tillenius et al. “Resource-Aware Task Scheduling”. In: *ACM Trans. Embed. Comput. Syst.* 14.1 (Jan. 2015), 5:1–5:25

After Mapping



Compilation Process



Task on process

- 1 Add declaration for each variable on each process
- 2 Add copy/communication for written variables
- 3 Substitute “original variables” by “local variables”
- 4 Remove “original variables” declarations

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```
int x;           int x;  
                 int x_0;  
                 ⇒ int x_1;  
                 int x_2;  
                 ...
```

Task on process

- 1 Add declaration for each variable on each process
- 2 **Add copy/communication for written variables**
- 3 Substitute “original variables” by “local variables”
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Copy/Communication

- inside the task
 - More precise
 - Issue for code generation on dynamic cases
- between the tasks
 - No dynamic cases
 - Less precise

Task on process

- 1 Add declaration for each variable on each process
- 2 Add copy/communication for written variables
- 3 **Substitute “original variables” by “local variables”**
- 4 Remove “original variables” declarations

```
#pragma distributed on_cluster 0
{
    ...
    x=0;
    ...
    x_0=x;
    x_1=x;
    x_2=x;
    ...
}

⇒

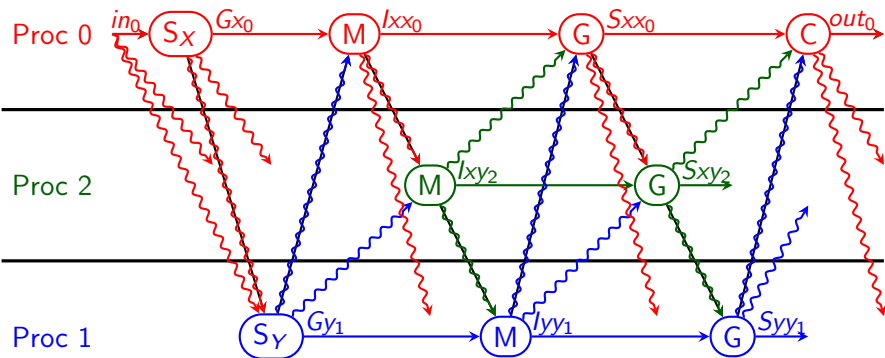
{
    ...
    x_0=0;
    ...
    x_0=x_0;
    x_1=x_0;
    x_2=x_0;
    ...
}
```

Task on process

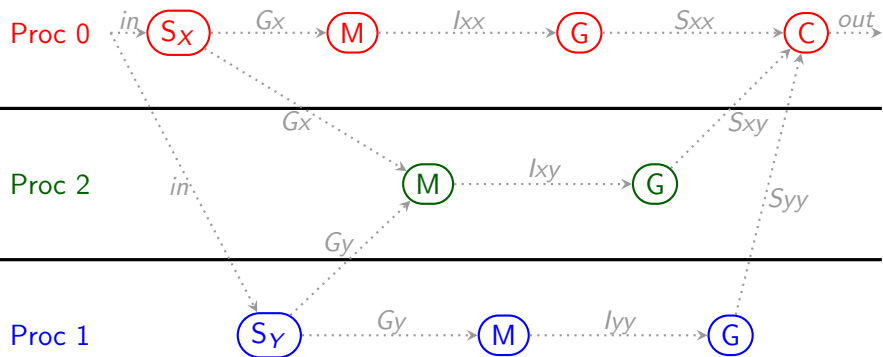
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int x_0;  
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int x_2;  
...  
⇒ int x_0;  
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   int x_2;  
   ...
```

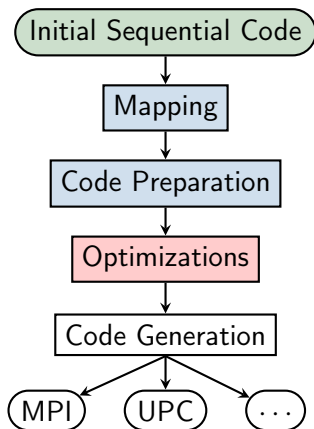
After Code Preparation



Task Graph



Compilation Process



Reduce Copy/Communication

- Dead-code Elimination
- Dead-iteration Elimination

Reduce Memory Footprint

- Array Resizing

Reduce Copy/Communication

- Dead-code Elimination
- Dead-iteration Elimination

	case 1	case 2	case 3
task on P	write x_P $x_Q = x_P$ $x_R = x_P$	write x_P $x_Q = x_P$ $x_R = x_P$	write $a_P[0..n]$ $a_Q[0..n] = a_P[0..n]$ $a_Q[0..n/2] = a_P[0..n/2]$ $a_R[0..n] = a_P[0..n]$ $a_R[n/2..n] = a_P[n/2..n]$
task on Q $Q \neq P$...	read x_Q write x_Q $x_P = x_Q$ $x_R = x_Q$	read $a_Q[0..n/2]$
task on R $R \neq \{P, Q\}$	read x_R	read x_R	read $a_R[n/2..n]$

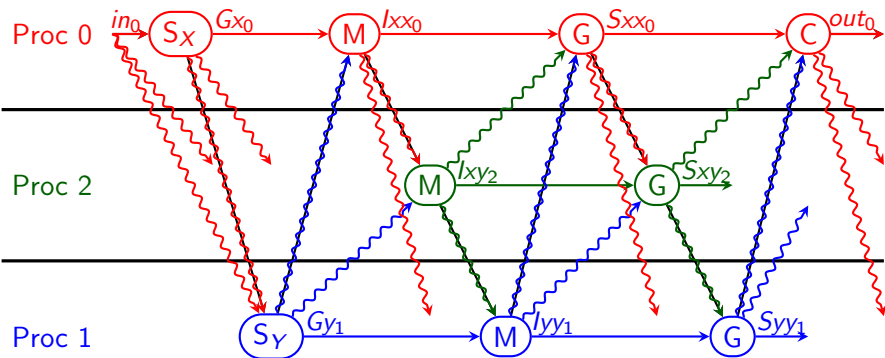
Reduce Memory Footprint

- Array Resizing
 - Compute new array size
 - Resize array declarations
 - Shift array cells access

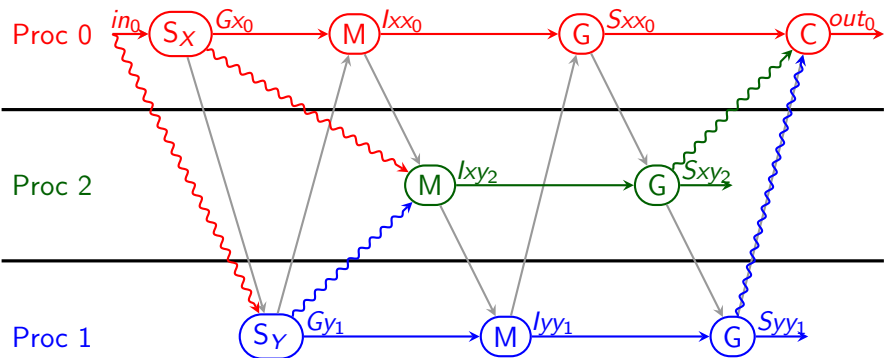
```
int a[20];  
int i;  
for (i=5; i<15; i++)  
    a[i] = i*i;
```

```
⇒ int a[10];  
   int i;  
   for (i=5; i<15; i++)  
       a[i-5] = i*i;
```

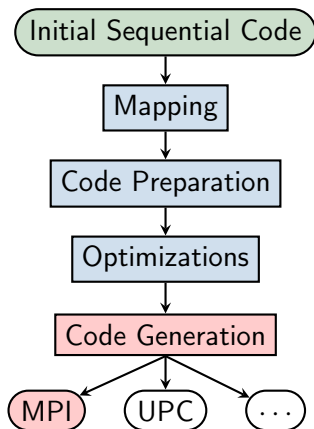
Before Optimizations



After Optimizations



Compilation Process



MPI Code Generation

- Configure MPI Environment
- Replace pragma block by test on process rank
- Replace copy by
 - Send message for rhs on rhs process to lhs process
 - Receive message for lhs on lhs process from rhs process

MPI Code Generation

- **Configure MPI Environment**
- Replace pragma block by test on process rank
- Replace copy by
 - Send message for rhs on rhs process to lhs process
 - Receive message for lhs on lhs process from rhs process

```
MPI_Status status;
int size, rank;
MPI_Init(&argc, &argv);
MPI_Comm_size(MPI_COMM_WORLD, &size);
MPI_Comm_rank(MPI_COMM_WORLD, &rank);
if (size < REQUIRE_PROC_NUMBER) {
    printf("Not enough processes launched!");
    MPI_Finalize();
    return 0;
}
...
MPI_Finalize();
return 0;
```

MPI Code Generation

- Configure MPI Environment
- **Replace pragma block by test on process rank**
- Replace copy by
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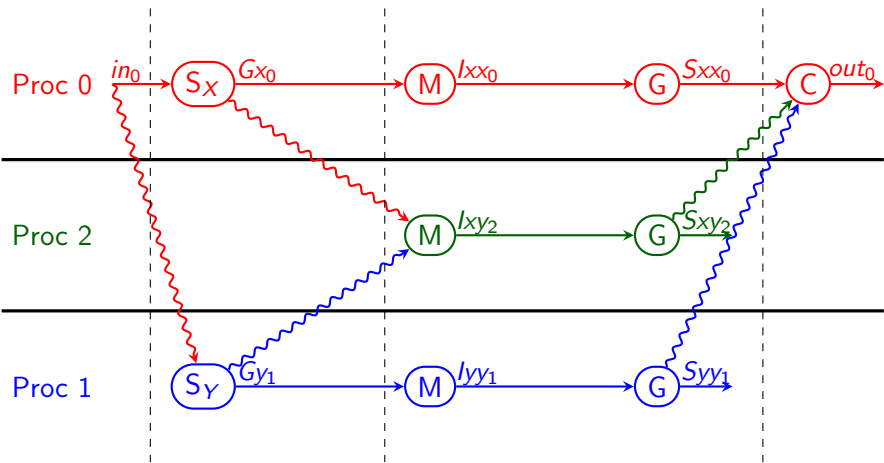
```
#pragma distributed on_cluster 0      if (rank==0)
{                                     ⇒ {
  ...                                 }
}
```

MPI Code Generation

- Configure MPI Environment
- Replace pragma block by test on process rank
- Replace copy by
 - **Send message for rhs on rhs process to lhs process**
 - **Receive message for lhs on lhs process from rhs process**

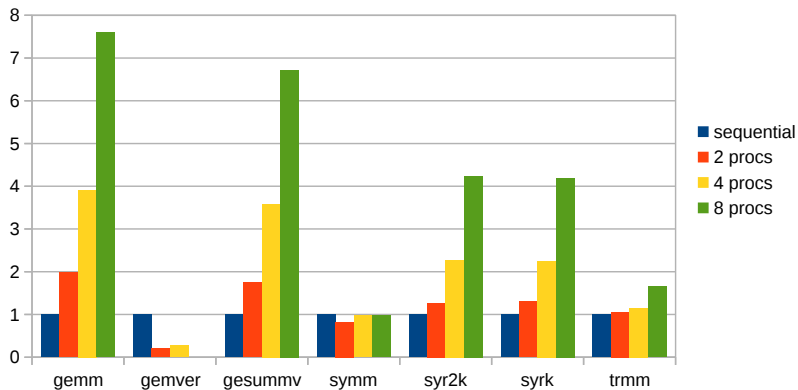
```
x_1 = x_0;           ⇒  
                      if (rank==0)  
                        MPI_Send(&x_0, 1, MPI_DOUBLE, 1, 0,  
                                MPI_COMM_WORLD);  
                      if (rank==1)  
                        MPI_Recv(&x_1, 1, MPI_DOUBLE, 0, 0,  
                                MPI_COMM_WORLD);
```

After Parallel Code Generation



Experimental Results

Speed up



Benchmark: BLAS in Polybench

size: $\sim 3000 \times 4000$

type: double

General Limitations

- Number of processes known at the beginning
- No dynamic parallelism
- Communication overestimation in case of dynamic communications

Experimental Limitations

- Strongly mapping dependent

Achievement

- Automatic source-to-source transformations
- Succession of simple transformations
- Basic communication functions
- Provable transformations
- Good efficiency

Future Work

- Improvement of the initial mapping by loop rescheduling
- Asynchronous communications instead of synchronous communications

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